

Outline

- Model checking of finite-state systems
- **Assisting in program development**
 - Program repair
 - Program differencing

Sound and Complete Mutation-Based Program Repair

Bat-Chen Rothenberg and Orna Grumberg

Formal Methods (FM'16)

Mutation-Based Program Repair

Sequential
program

Assertions
in code

Given set
of
mutations

Can we use
these
mutations to
make all
assertions
hold?

Assignments,
conditionals,
loops and
function calls



Assertion
violation

operator
replacement
($+$ \rightarrow $-$),
constant
manipulation
($c \rightarrow c + 1$)

Return
all
possible
repairs

Example

```
int f(int x, int y){  
1.   int z;  
2.   if (x + y > 8) {  
3.       z = x + y;  
4.   } else {  
5.       z = 9;  
6.   }  
7.   if (z ≥ 9) z = z - 1;  
8.   assert(z > 8);  
9.   return z;  
   }
```

$x = 5, y = 2$

$z = 9$

$z = 8$



Example

```
int f(int x, int y){  
1.   int z;  
2.   if (x + y > 8) {  
3.       z = x + y;  
4.   } else {  
5.       z = 9;  
6.   }  
7.   if (z ≥ 9) z = z + 1;  
8.   assert(z > 8);  
9.   return z;  
}
```

At
this
point
 $z \geq 9$

Mutation list:

Replace + with -
Replace - with +
Replace > with ≥
Replace ≥ with >

Repair list:

option 1:

line 7: replace ≥ with >

option 2:

line 7: replace - with +

Note:
Repairs
are
minimal



Example

```
int f(int x, int y){  
1.   int z;  
2.   if (x + y > 9) {  
3.       z = x + y;  
4.   } else {  
5.       z = 10;  
6.   }  
7.   if (z ≥ 9) z = z - 1;  
8.   assert(z > 8);  
9.   return z;  
   }
```

At this
point
 $z \geq 10$



Mutation list:

Replace + with -

Replace - with +

Replace > with \geq

Replace \geq with >

Increase constants by 1

Program Repair

- **Programs:** sequential C programs
- **Specification:** Assertions added in program text
- **Bug:** A program run which violates an assertion
- **Repair:** Changed statement(s) in the program, resulting in a correct program

- **No assumption** on the number of faulty expressions.

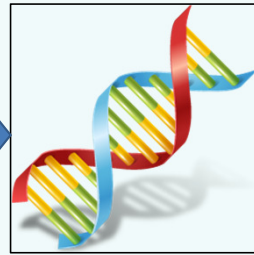
- **Goal:** To return a sequence of repaired (correct) programs in increasing number of changes using **techniques and tools of formal methods**

Overview of our approach

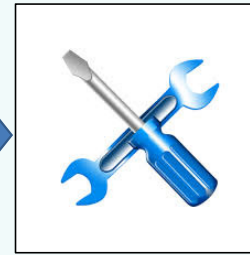
```
int f(int x, int y){  
1.   int z;  
2.   if (x + y > 8) {  
3.       z = x + y;  
4.   } else {  
5.       z = 9;  
6.   }  
7.   if (z ≥ 9) z = z - 1;  
8.   assert(z > 8);  
9.   return z;  
}
```



Translation



Mutation



Repair

line 7: replace operator \geq with $>$
line 7: replace operator $-$ with $+$
...

Output:

All **minimal** repairs,
sorted by size

Input:
a buggy
program

Finding all **unsatisfiable constraint sets**
from a finite set of **programs**



First step - Translation

Goal: Translate the program into a **set of constraints** which is **satisfiable iff the program has a bug** (i.e. there exists an input for which an assertion fails)

Work by Clarke, Kroening, Lerda (TACAS 2000)
(CBMC)

- Simplification
- Unwinding of loops
 - a **bounded** number of unwinding
- Conversion to SSA

Correctness
is bounded

First step - Translation

```
int f(int x, int y){  
1.   int z;  
2.   if (x + y > 8) {  
3.       z = x + y;  
4.   } else {  
5.       z = 9;  
6.   }  
7.   if (z ≥ 9) z = z - 1;  
8.   assert(z > 8);  
9.   return z;  
}
```



```
{ g1 = x1 + y1 > 8,  
  z2 = x1 + y1,  
  z3 = 9,  
  z4 = g1?z2:z3,  
  b1 = z4 ≥ 9 ,  
  z5 = z4 - 1,  
  z6 = b1?z5:z4,  
  z6 ≤ 8  
}
```

[Clarke et al. 2004 (CBMC)]

First step - Translation

```
int f(int x, int y){  
1.   int z;  
2.   if (x + y > 8) {  
3.       z = x + y;  
4.   } else {  
5.       z = 9;  
6.   }  
7.   if (z ≥ 9) z = z - 1;  
8.   assert(z > 8);  
9.   return z;  
}
```



```
{  $g_1 = x_1 + y_1 > 8,$   
   $z_2 = x_1 + y_1,$   
   $z_3 = 9,$   
   $z_4 = g_1 ? z_2 : z_3,$   
   $b_1 = z_4 \geq 9,$   
   $z_5 = z_4 - 1,$   
   $z_6 = b_1 ? z_5 : z_4,$   
   $z_6 \leq 8$   
}
```

[Clarke et al. 2004 (CBMC)]

First step - Translation

```
int f(int x, int y){  
1.   int z;  
2.   if (x + y > 8) {  
3.       z = x + y;  
4.   } else {  
5.       z = 9;  
6.   }  
7.   if (z ≥ 9) z = z - 1;  
8.   assert(z > 8);  
9.   return z;  
}
```



```
{ g1 = x1 + y1 > 8,  
  z2 = x1 + y1,  
  z3 = 9,  
  z4 = g1?z2:z3,  
  b1 = z4 ≥ 9,  
  z5 = z4 - 1,  
  z6 = b1?z5:z4,  
  z6 ≤ 8  
}
```

[Clarke et al. 2004 (CBMC)]

First step - Translation

```
int f(int x, int y){  
1.   int z;  
2.   if (x + y > 8) {  
3.       z = x + y;  
4.   } else {  
5.       z = 9;  
6.   }  
7.   if (z ≥ 9) z = z - 1;  
8.   assert(z > 8);  
9.   return z;  
}
```



```
{ g1 = x1 + y1 > 8,  
  z2 = x1 + y1,  
  z3 = 9,  
  z4 = g1?z2:z3,  
  b1 = z4 ≥ 9,  
  z5 = z4 - 1,  
  z6 = b1?z5:z4,  
  z6 ≤ 8  
}
```

[Clarke et al. 2004 (CBMC)]

First step - Translation

```
int f(int x, int y){  
1.   int z;  
2.   if (x + y > 8) {  
3.       z = x + y;  
4.   } else {  
5.       z = 9;  
6.   }  
7.   if (z ≥ 9) z = z - 1;  
8.   assert(z > 8);  
9.   return z;  
}
```



```
{ g1 = x1 + y1 > 8,  
  z2 = x1 + y1,  
  z3 = 9,  
  z4 = g1? z2: z3,  
  b1 = z4 ≥ 9 ,  
  z5 = z4 - 1,  
  z6 = b1? z5: z4,  
  z6 ≤ 8  
}
```

[Clarke et al. 2004 (CBMC)]



```
int f(int x, int y){
1.   int z;
2.   if (x + y > 8) {
3.       z = x + y;
4.   } else {
5.       z = 9;
6.   }
7.   if (z ≥ 9) z = z - 1;
8.   assert(z > 8);
9.   return z;
}
```

simplification

```
int f(int x, int y){
1.   int z;
2.1  bool g = x + y > 8;
2.2  if (g) {
3.       z = x + y;
4.   } else {
5.       z = 9;
6.   }
7.1  bool b = z ≥ 9
7.2  if (b) z = z - 1;
8.   assert(z > 8);
9.   return z;
}
```

Conversion
to SSA

```
int f(int x1, int y1){
1.   int z1;
2.1  bool g1 = x1 + y1 > 8;
2.2  if (g1) {
3.       z2 = x1 + y1;
4.   } else {
5.       z3 = 9;
6.   }
-   z4 = g1?z2:z3;
7.1  bool b1 = z4 ≥ 9
7.2  if (b1) z5 = z4 - 1;
-   z6 = b1?z5:z4;
8.   assert(z6 > 8);
9.   return z6;
}
```

```
{ g1 = x1 + y1 > 8,
  z2 = x1 + y1,
  z3 = 9,
  z4 = g1?z2:z3,
  b1 = z4 ≥ 9,
  z5 = z4 - 1,
  z6 = b1?z5:z4,
  z6 ≤ 8
}
```

What about loops?

```
int z = 0;
while(x > 0){
  z = z + x;
  x = x - 1;
}
```

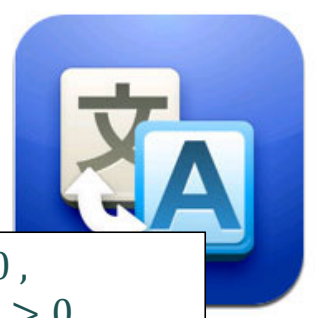
simplification

```
int z = 0;
bool t = x > 0;
while(t){
  z = z + x;
  x = x - 1;
  t = x > 0;
}
```

unwinding
(for $b = 2$)

```
int z = 0;
bool t = x > 0;
if(t){
  z = z + x;
  x = x - 1;
  t = x > 0;
  if(t){
    z = z + x;
    x = x - 1;
    t = x > 0;
    assume(!t);
  }
}
```

```
{ z1 = 0,
  t1 = x1 > 0,
  z2 = z1 + x1},
  x2 = x1 - 1,
  t2 = x2 > 0,
  z3 = z2 + x2},
  x3 = x2 - 1,
  t3 = x3 > 0,
  t1 ∧ t2 → !t3},
  z4 = t2?z3:z2},
  x4 = t2?x3:x2},
  z5 = t1?z4:z1},
  x5 = t1?x4:x1}
}
```



Translation

- In the translation, loops are unwound a **bounded** number of times
- **Important observation: correctness is bounded.**
That is, repairs found by our method only guarantee that assertions cannot be violated by inputs going through the loop at most **k times**

More complex example

```
int f(int x, int y){  
1.     int z;  
2.     if (x + y > 8) {  
3.         z = x + y;  
4.     } else {  
5.         z = 9;  
6.     }  
7.     while(x > 0){  
8.         z = z + x;  
9.         x = x - 1;  
10.    }  
11.    if (z ≥ 9) z = z - 1;  
12.    assert(z > 8);  
13.    return z;  
}
```



```
{ g1 = x1 + y1 > 8,  
  z2 = x1 + y1,  
  z3 = 9,  
  z4 = g1?z2:z3,  
  t1 = x1 > 0,  
  z5 = z4 + x1,  
  x2 = x1 - 1,  
  t2 = x2 > 0,  
  z6 = z5 + x2,  
  x3 = x2 - 1,  
  t3 = x3 > 0,  
  t1 ∧ t2 → !t3,  
  z7 = t2?z6:z5,  
  z8 = t1?z7:z4,  
  b1 = z8 ≥ 9,  
  z9 = z8 - 1,  
  z10 = b1?z9:z8,  
  z10 ≤ 8  
}
```

Let's see how mutations to the program affect the set of constraints

```

int f(int x, int y){
1.   int z;
2.   if (x + y > 8) {
3.       z = x + y;
4.   } else {
5.       z = 9;
6.   }
7.   while(x > 0){
8.       z = z - x;
9.       x = x - 1;
10.  }
11.  if (z ≥ 9) z = z - 1;
12.  assert(z > 8);
13.  return z;
}

```

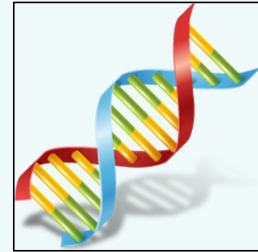


```

{ g1 = x1 + y1 > 8,
  z2 = x1 + y1,
  z3 = 9,
  z4 = g1?z2:z3,
  t1 = x1 > 0,
  z5 = z4 - x1,
  x2 = x1 - 1,
  t2 = x2 > 0,
  z6 = z5 - x2,
  x3 = x2 - 1,
  t3 = x3 > 0,
  t1 ∧ t2 → !t3,
  z7 = t2?z6:z5,
  z8 = t1?z7:z4,
  b1 = z8 ≥ 9,
  z9 = z8 - 1,
  z10 = b1?z9:z8,
  z10 ≤ 8
}

```

Second step- Mutation

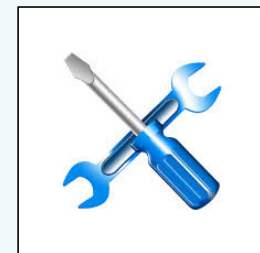


$$\{ \begin{array}{l} g_1 = x_1 + y_1 > 8, \\ z_2 = x_1 + y_1, \\ z_3 = 9, \\ z_4 = g_1 ? z_2 : z_3, \\ t_1 = x_1 > 0, \\ z_5 = z_4 + x_1, \\ x_2 = x_1 - 1, \\ t_2 = x_2 > 0, \\ z_6 = z_5 + x_2, \\ x_3 = x_2 - 1, \\ t_3 = x_3 > 0, \\ t_1 \wedge t_2 \rightarrow !t_3, \\ z_7 = t_2 ? z_6 : z_5, \\ z_8 = t_1 ? z_7 : z_4, \\ b_1 = z_8 \geq 9, \\ z_9 = z_8 - 1, \\ z_{10} = b_1 ? z_9 : z_8, \\ z_{10} \leq 8 \end{array} \}$$


S_{soft}

$$\{ \begin{array}{l} g_1 = x_1 + y_1 > 8, \\ z_2 = x_1 + y_1, \\ z_3 = 9, \\ t_1 = x_1 > 0, \\ z_5 = z_4 + x_1, \\ x_2 = x_1 - 1, \\ t_2 = x_2 > 0, \\ z_6 = z_5 + x_2, \\ x_3 = x_2 - 1, \\ t_3 = x_3 > 0, \\ b_1 = z_8 \geq 9, \\ z_9 = z_8 - 1 \end{array} \}$$


$$\{ \begin{array}{l} g_1 = x_1 + y_1 > 8, \\ z_2 = x_1 + y_1, \\ z_3 = 9, \\ t_1 = x_1 > 0 \wedge t_2 = x_2 > 0 \wedge t_3 = x_3 > 0, \\ z_5 = z_4 + x_1 \wedge z_6 = z_5 + x_2, \\ x_2 = x_1 - 1 \wedge x_3 = x_2 - 1, \\ b_1 = z_8 \geq 9, \\ z_9 = z_8 - 1 \end{array} \}$$

$$\{ \begin{array}{l} z_4 = g_1 ? z_2 : z_3, \\ t_1 \wedge t_2 \rightarrow !t_3, \\ z_7 = t_2 ? z_6 : z_5, \\ z_8 = t_1 ? z_7 : z_4, \\ z_{10} = b_1 ? z_9 : z_8, \\ z_{10} \leq 8 \end{array} \}$$


$$g_1 = x_1 + y_1 > 8$$

Second step - Mutation

Mutation list:
 Replace + with -
 Replace - with +
 Replace > with ≥
 Replace ≥ with >

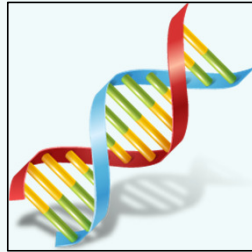
```

int f(int x, int y){
1.   int z;
2.   if (x + y > 8) {
3.       z = x - y;
4.   } else {
5.       z = 9;
6.   }
7.   if (z ≥ 9) {
           z = z - 1;
           }
8.   assert(z > 8);
9.   return z;
}

```

$\{g_1 = x_1 + y_1 > 8, g_1 = x_1 + y_1 \geq 8, g_1 = x_1 + y_1 > 8, g_1 = x_1 + y_1 \geq 8\}$
 $\{z_2 = x_1 - y_1, z_2 = x_1 - y_1\}$
 $\{z_3 = 9\}$
 $z_4 = g_1 ? z_2 : z_3$
 $\{b_1 = z_4 \geq 9, b_1 = z_4 > 9\}$
 $\{z_5 = z_4 - 1, z_5 = z_4 + 1\}$
 $z_6 = b_1 ? z_5 : z_4$
 $z_6 \leq 8$

Second step- Mutation



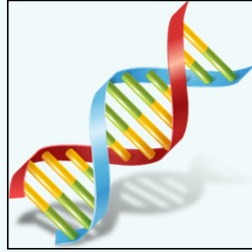
$$\begin{cases}
 g_1 = x_1 + y_1 > 8, \\
 z_2 = x_1 + y_1, \\
 z_3 = 9, \\
 \mathbf{t_1 = x_1 > 0 \wedge t_2 = x_2 > 0 \wedge t_3 = x_3 > 0}, \\
 z_5 = z_4 + x_1 \wedge z_6 = z_5 + x_2, \\
 x_2 = x_1 - 1 \wedge x_3 = x_2 - 1, \\
 b_1 = z_8 \geq 9, \\
 z_9 = z_8 - 1 \\
 \}
 \end{cases}$$



$$\begin{aligned}
 S_1 &= \{ g_1 = x_1 + y_1 > 8, g_1 = x_1 - y_1 > 8, \\
 &\quad g_1 = x_1 + y_1 \geq 8 \} \\
 S_2 &= \{ z_2 = x_1 + y_1, z_2 = x_1 - y_1 \} \\
 S_3 &= \{ z_3 = 9 \} \\
 S_4 &= \{ \mathbf{t_1 = x_1 > 0 \wedge t_2 = x_2 > 0 \wedge t_3 = x_3 > 0}, \\
 &\quad \mathbf{t_1 = x_1 \geq 0 \wedge t_2 = x_2 \geq 0 \wedge t_3 = x_3 \geq 0} \} \\
 S_5 &= \{ z_5 = z_4 + x_1 \wedge z_6 = z_5 + x_2, \\
 &\quad z_5 = z_4 - x_1 \wedge z_6 = z_5 - x_2 \} \\
 S_6 &= \{ x_2 = x_1 - 1 \wedge x_3 = x_2 - 1, \\
 &\quad x_2 = x_1 + 1 \wedge x_3 = x_2 + 1 \} \\
 S_7 &= \{ b_1 = z_8 \geq 9, b_1 = z_8 > 9 \} \\
 S_8 &= \{ z_9 = z_8 - 1, z_9 = z_8 + 1 \}
 \end{aligned}$$

Replace + with -
 Replace - with +
 Replace > with \geq
 Replace \geq with >

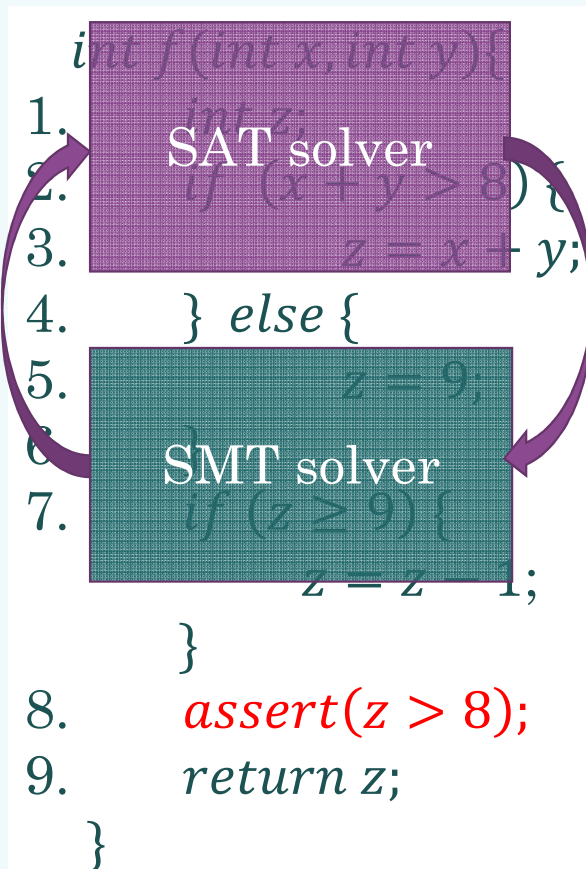
Second step- Mutation



We have **reduced** the problem of **finding a correct program** from a finite set of mutated programs to the problem of **choosing one constraint from each S_i** such that the conjunction of all chosen constraints and all constraints in S_{hard} is **unsatisfiable**.



Third step - Repair



$\{g_1 = x_1 + y_1 > 8, g_1 = x_1 - y_1 > 8, g_1 = x_1 + y_1 \geq 8\}$

$\{z_2 = x_1 + y_1, z_2 = x_1 - y_1\}$

$\{z_3 = 9\}$

$z_4 = g_1 ? z_2 : z_3$

$\{b_1 = z_4 \geq 9, b_1 = z_4 > 9\}$

$\{z_5 = z_4 - 1, z_5 = z_4 + 1\}$

$z_6 = b_1 ? z_5 : z_4$

$z_6 \leq 8$



Repair

SAT solver

Choose candidate program of **size = 1**
 Blocking clause for specific assignment
 Blocking clause for this assignment
 And all other supersets of changes

SMT solver

$g_1 = x_1 + y_1 > 8$
 $z_4 = g_1 ? z_2 : z_3$ $z_2 = x_1 + y_1$
 $z_6 = b_1 ? z_5 : z_4$ $z_3 = 9$
 $z_6 \leq 8$ $b_1 = z_4 \geq 9$
 $z_5 = z_4 - 1$

SAT
 $c_1 = 0$
 $c_2 = 0$
 $c_3 = 0$
 $c_4 = 1$
 $c_5 = 0$
 $c_6 = 1$
 $c_7 = 0$
 $c_8 = 0$
 $c_9 = 1$
 $c_{10} = 0$

c_1 c_2
 $\{g_1 = x_1 + y_1 > 8, g_1 = x_1 - y_1 > 8,$
 $g_1 = x_1 + y_1 \geq 8\}$
 c_3

c_4 c_5
 $\{z_2 = x_1 + y_1, z_2 = x_1 - y_1\}$ UNSAT

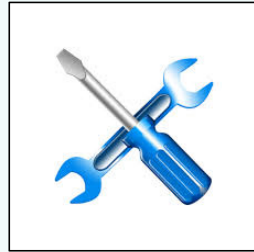
c_6
 $\{z_3 = 9\}$

c_7 c_8
 $\{b_1 = z_4 \geq 9, b_1 = z_4 > 9\}$

c_9 c_{10}
 $\{z_5 = z_4 - 1, z_5 = z_4 + 1\}$

UNSAT
 ((repair
 failed!))

Final step- Repair



$S_1 = \{1, 2, 3\}$
 $S_2 = \{4, 5\}$
 $S_3 = \{6\}$
 $S_4 = \{7, 8\}$
 $S_5 = \{9, 10\}$
 $S_6 = \{11, 12\}$
 $S_7 = \{13, 14\}$
 $S_8 = \{15, 16\}$

- Each set S_i contains a **constraint** c_o^i , encoding the original (unmutated) statement (marked in **red** in the example)
- A **selection vector (SV)** is a vector of constraints $[c_1, \dots, c_n]$ where c_i is taken from S_i for all $1 \leq i \leq n$. For example: $v' = [2, 4, 6, 7, 9, 12, 13, 15]$
- **Note that each SV “encodes” a program**
- a SV is said to be **correct** if it encodes a (bounded) correct program.

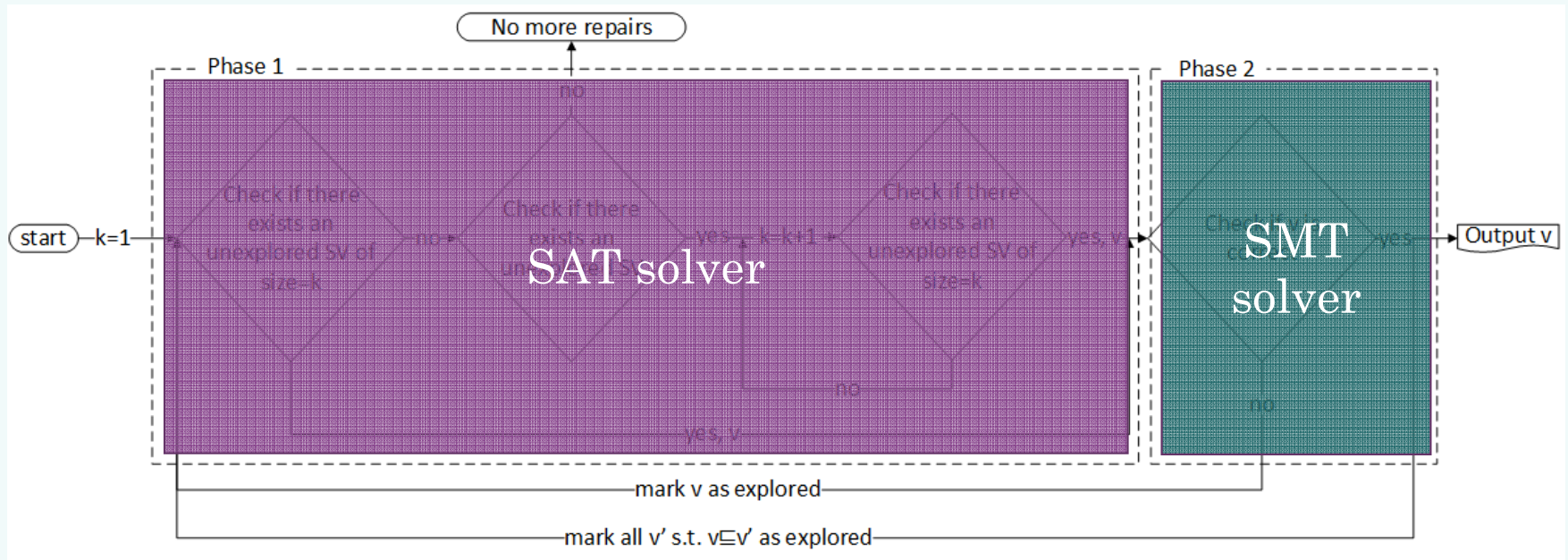
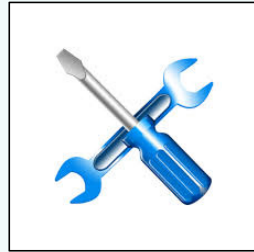
Final step- Repair



$$\begin{aligned} S_1 &= \{1, 2, 3\} \\ S_2 &= \{4, 5\} \\ S_3 &= \{6\} \\ S_4 &= \{7, 8\} \\ S_5 &= \{9, 10\} \\ S_6 &= \{11, 12\} \\ S_7 &= \{13, 14\} \\ S_8 &= \{15, 16\} \end{aligned}$$

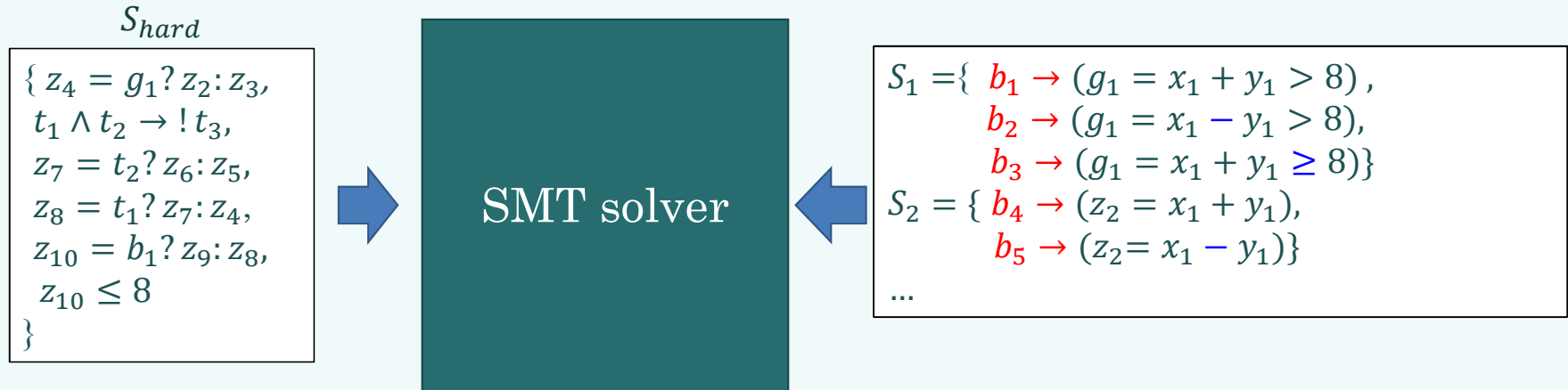
- The **size** of a selection vector v is the number of sets from which the chosen constraint is different than c_0^i , i.e. the number of lines changed in the program encoded by v . For example: $\text{size}(v') = 2$
- Size order: $[2, 4, 6, 7, 9, 12, 13, 15] \sqsubseteq [2, 4, 6, 7, 9, 12, 14, 15]$
but $[2, 4, 6, 7, 9, 12, 13, 15] \not\sqsubseteq [3, 4, 6, 7, 9, 12, 14, 15]$
- A SV v is a **Minimal correct SV (MCSV)** if v is a correct SV and there does not exist a SV v' s.t. $v' \sqsubseteq v$.
- **We would like to return all minimal SVs in increasing size order.**

Repair scheme

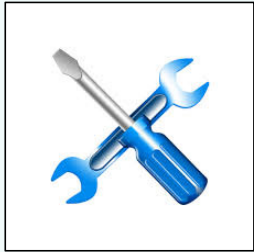




Initialization of the SMT solver



The formula solved by the SAT solver will be over the Boolean variables we added.



Initialization of the SAT solver

Every satisfying assignment corresponds to a SV

$$\begin{aligned} b_1 + b_2 + b_3 = 1 \wedge \\ b_4 + b_5 = 1 \wedge \\ \dots \end{aligned}$$



All SV's returned are of size $\leq k$

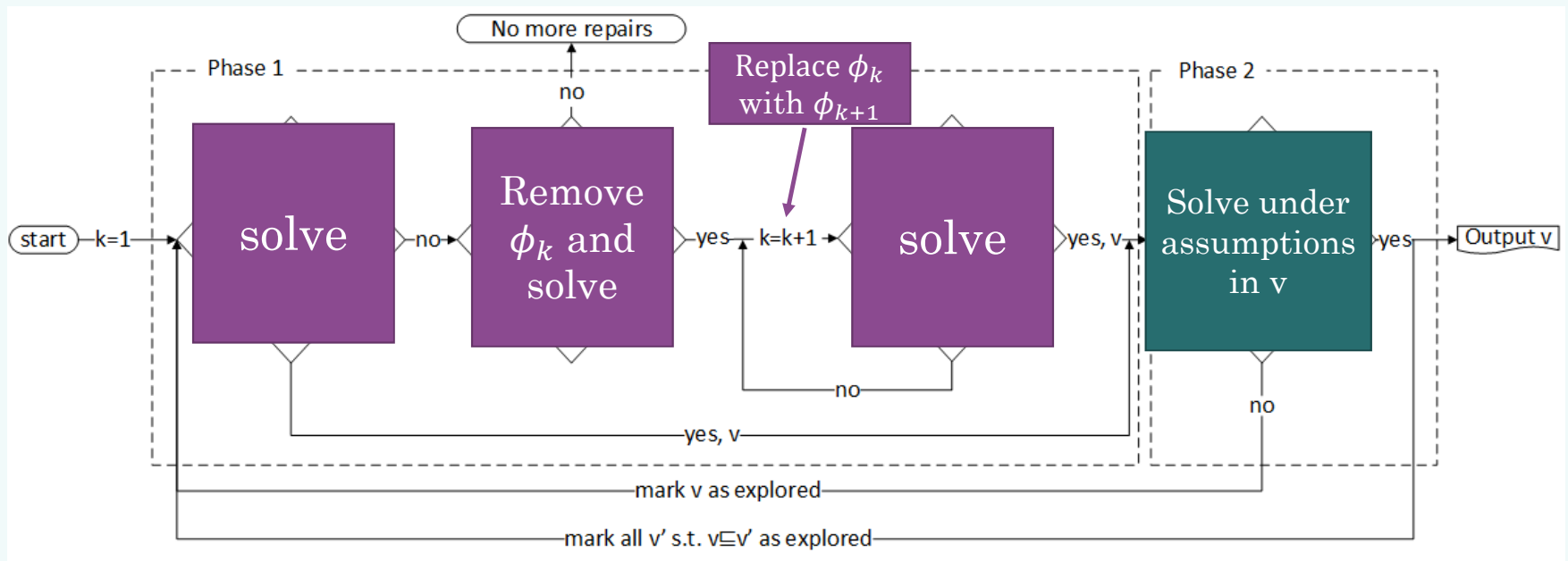
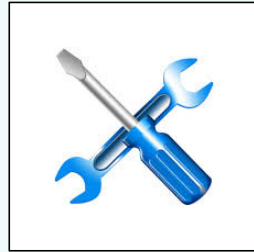
$$b_1 + b_4 + \dots \geq n - k$$



This constraint, called **cardinality constraint**, will be referred to as ϕ_k

$$\begin{aligned} S_1 = \{ & b_1 \rightarrow (g_1 = x_1 + y_1 > 8), \\ & b_2 \rightarrow (g_1 = x_1 - y_1 > 8), \\ & b_3 \rightarrow (g_1 = x_1 + y_1 \geq 8) \} \\ S_2 = \{ & b_4 \rightarrow (z_2 = x_1 + y_1), \\ & b_5 \rightarrow (z_2 = x_1 - y_1) \} \\ & \dots \end{aligned}$$

Repair scheme



Results on our example

Program

```
int f(int x, int y){  
1.   int z;  
2.   if (x + y > 8) {  
3.       z = x + y;  
4.   } else {  
5.       z = 9;  
6.   }  
7.   while(x > 0){  
8.       z = z + x;  
9.       x = x - 1;  
10.  }  
11.  if (z ≥ 9) z = z - 1;  
12.  assert(z > 8);  
13.  return z;  
}
```

List of mutations

Replace + with -
Replace - with +
Replace > with ≥
Replace ≥ with >

Unwinding bound (not including)

3

Results

line 11: replace operator ≥ with >
line 11: replace - with +

Results on our example

Program

```
int f(int x, int y){
1.     int z;
2.     if (x + y  $\geq$  8) {
3.         z = x - y;
4.     } else {
5.         z = 9;
6.     }
7.     while(x > 0){
8.         z = z + x;
9.         x = x - 1;
10.    }
11.    if (z  $\geq$  9) z = z - 1;
12.    assert(z > 8);
13.    return z;
}
```

List of mutations

Replace + with -
Replace - with +
Replace > with \geq
Replace \geq with >



Results

line 2: replace \geq with > *and*
line 3 replace - with + *and*
line 11 replace - with +
line 2: replace \geq with > *and*
line 3 replace - with + *and*
line 11 replace \geq with >

Unwinding bound (not including)

3

Theorem:

Our algorithm is sound and complete.

That is, for a given bound \mathbf{b} :

A program is returned by our algorithm
iff

it is **minimal** and **\mathbf{b} -bounded correct**

- Minimal number of changes
- Every assertion reachable along a computation of bounded length \mathbf{b} is correct

Ver.	Method of [11]		Method of [12]		Our method			
	Fixed?	Time[s]	Fixed?	Time[s]	Fixed?	Time[s]	Fixed?	Time[s]

3									
6									
7									
8									
9									
10									
12									
16									
17									
18									
19									
20									
25									
28	+	34	+	35			+	93.678	
31					+	1.246	+	4.661	
32					+	1.902	+	85.349	
35	+	41	+	46			+	92.866	
36	+	8	+	6			+	94.599	
39	+	82	+	101	+	2.558	+	16.393	
40									

		Level 1	Level 2
Op. replacement	Arithmetic	{+, -}, {*, /, %}	{+, -, *, /, %}
	Relational	{>, >=}, {<, <=}	{>, >=, <, <=}, {==, !=}
	Logical	{ , &&}	
	Bit-wise	{>>, <<}, {&, , ^}	
Constant manipulation			C → C+1, C → C-1, C → -C, C → 0

implemented in the tool FoREnSiC.

16 (39%)	38	15 (36.6%)	38	11 (26.83%)	2.278	18 (43.9%)	48.151
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Summary

- We suggest a repair method which returns **all minimal** (bounded) correct programs, **in increasing order**
 - Based on a given set of mutations
- If no repaired program is returned then the given mutations **cannot repair** the program
- SAT solver handles the search of mutated programs
- SMT solver checks if a mutated program is correct
- Both solvers are used **incrementally**

Summary

- Minimal mutations: No change is made to the original program unless necessary
- The method can assist a programmer in debugging in initial stages of development
 - When bugs are simple, but many

Questions?